

The Generic Chaining

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ISBN 3-540-24518-9

The Appendix was not published in the book.

Therefore, the Appendix is shown on the next pages.

A. Appendix: Trees and Majorizing Measures

In this appendix we describe different ways to measure the size of a metric space. Some of these ways played an important part in the development of the theory. We will show that they are all equivalent to the functional $\gamma_2(T, d)$. It is possible to consider more general notions corresponding to other functionals considered in the book, but for simplicity we consider only the case of γ_2 .

A *tree* \mathcal{T} of a metric space (T, d) is a *finite* collection of subsets of T with the following two properties.

Given A, B in \mathcal{T} , if $A \cap B \neq \emptyset$, then either $A \subset B$ or else $B \subset A$. (A.1)

\mathcal{T} has a largest element . (A.2)

If $A, B \in \mathcal{T}$ and $B \subset A$, $B \neq A$, we say that B is a *child* of A if

$$C \in \mathcal{T}, B \subset C \subset A \Rightarrow C = B \text{ or } C = A. \quad (\text{A.3})$$

We denote by $c(A)$ the number of children of A . We will consider only trees with the following property

If $A \in \mathcal{T}$ and $c(A) = 0$, then A contains exactly one point . (A.4)

A *separated* tree is a tree \mathcal{T} such that to each A in \mathcal{T} with $c(A) \geq 1$ is associated an integer $s(A) \in \mathbb{Z}$ with the following properties.

If B_1 and B_2 are children of A , then $d(B_1, B_2) \geq 4^{-s(A)}$. (A.5)

If B is a child of A , then $s(B) > s(A)$. (A.6)

Here of course $d(B_1, B_2) = \inf\{d(x_1, x_2), x_1 \in B_1, x_2 \in B_2\}$. We then define

$$S_{\mathcal{T}} = \{t \in T ; \{t\} \in \mathcal{T}\}$$

and the *depth* of \mathcal{T} ,

$$d(\mathcal{T}) = \inf_{t \in S_{\mathcal{T}}} \sum_{t \in A \in \mathcal{T}} 4^{-s(A)} \sqrt{\log c(A)} .$$

Here and below we make the convention that the summation does not include the term $A = \{t\}$ (for which $c(A) = 0$). We view $d(\mathcal{T})$ as a “measure

of the size" of the separated tree \mathcal{T} . We can then measure the size of T by $\sup\{d(\mathcal{T}) ; \mathcal{T} \text{ separated tree}\}$.

The notion of tree we just considered is but one of many possible. Let us now consider another (more restrictive) notion. An *organized* tree is a tree \mathcal{T} such that to each $A \in \mathcal{T}$ with $c(A) \geq 1$ are associated $j = j(A) \in \mathbb{Z}$, $t \in T$ and $t_1, \dots, t_{c(A)} \in B(t, 4^{-j})$ with the properties that

$$1 \leq \ell < \ell' \leq c(A) \Rightarrow d(t_\ell, t_{\ell'}) \geq 4^{-j-1}$$

and that each ball $B(t_\ell, 4^{-j-2})$ contains exactly one child of A .

If B_1 and B_2 are children of A , then

$$d(B_1, B_2) \geq 4^{-j(A)-2}, \quad (\text{A.7})$$

so that an organized tree is also a separated tree (with $s(A) = j(A) + 2$), but the notion of organized tree is more restrictive. (For example we have no control of the diameter of the children of A in a separated tree.)

We define the depth $d'(\mathcal{T})$ of an organized tree by

$$d'(\mathcal{T}) = \inf_{t \in S_{\mathcal{T}}} \sum_{t \in A \in \mathcal{T}} 4^{-j(A)} \sqrt{\log c(A)}.$$

If we simply view \mathcal{T} as a separated tree using (A.7), then $d(\mathcal{T}) = d'(\mathcal{T})/16$ (where $d(\mathcal{T})$ is the depth of \mathcal{T} as a separated tree). Thus we have shown the following.

Proposition A.1. *We have*

$$\sup\{d'(\mathcal{T}) ; \mathcal{T} \text{ organized tree}\} \leq 16 \sup\{d(\mathcal{T}) ; \mathcal{T} \text{ separated tree}\}. \quad (\text{A.8})$$

Proposition A.2. *We have*

$$\gamma_2(T, d) \leq L \sup\{d'(\mathcal{T}) ; \mathcal{T} \text{ organized tree}\}. \quad (\text{A.9})$$

Proof. We consider the functional

$$F_n(A) = F(A) = \sup\{d'(\mathcal{T}) ; \mathcal{T} \subset A, \mathcal{T} \text{ organized tree}\},$$

where we write $\mathcal{T} \subset A$ as a shorthand for " $\forall B \in \mathcal{T}, B \subset A$ ". In the course of the proof of Theorem 1.3.1 we have noted that this theorem holds true as soon as (1.31) holds true when a is of the type r^{-j-1} . We check this condition when $r = 4$, $\theta(n) = 2^{n/2-2}$, $\beta = 1$, and $\tau = 1$. Consider $n \geq 0$ and $m = N_{n+1}$. Consider $j \in \mathbb{Z}$, $t \in T$ and $t_1, \dots, t_m \in B(t, 4^{-j})$ with

$$1 \leq \ell < \ell' \leq m \Rightarrow d(t_\ell, t_{\ell'}) \geq 4^{-j-1}.$$

Consider sets $H_\ell \subset B(t_\ell, 4^{-j-2})$ and $c < \min_{\ell \leq m} F(H_\ell)$. Consider, for $\ell \leq m$ a tree $\mathcal{T}_\ell \subset H_\ell$ with $d'(\mathcal{T}_\ell) > c$ and denote by A_ℓ its largest element. Then it should be obvious that the tree \mathcal{T} consisting of $U = \bigcup_{\ell \leq m} H_\ell$ (its

largest element) and the reunion of the trees \mathcal{T}_ℓ , $\ell \leq m$, is organized (with $j(U) = j$, and A_1, \dots, A_m as children of U). Moreover $S_{\mathcal{T}} = \bigcup_{\ell \leq m} S_{\mathcal{T}_\ell}$.

Consider $t \in S_{\mathcal{T}}$, and let ℓ with $t \in S_{\mathcal{T}_\ell}$. Then

$$\begin{aligned} \sum_{t \in A \in \mathcal{T}} 4^{-j(A)} \sqrt{\log c(A)} &= 4^{-j} \sqrt{\log m} + \sum_{t \in A \in \mathcal{T}_\ell} 4^{-j(A)} \sqrt{\log c(A)} \\ &\geq 4^{-j} \sqrt{\log m} + d'(\mathcal{T}_\ell) \geq 4^{-j} \sqrt{\log m} + c. \end{aligned}$$

Since $\sqrt{\log m} \geq 2^{n/2}$, this proves (1.31).

To prove (A.9) we apply Lemma 1.3.3. To control the diameter of T , we simply note that if $s, t \in T$, and j is the largest integer with $4^{-j} \geq d(s, t)$, then the tree \mathcal{T} consisting of $T, \{t\}, \{s\}$, is organized with $j = j(T)$ and $c(T) = 2$, so $d'(\mathcal{T}) \geq 4^{-j} \sqrt{\log 2}$. \square

Proposition A.3. *Given a metric space (T, d) we can find on T a probability measure μ , supported by a countable subset of T , and such that*

$$\sup_{t \in T} \int_0^\infty \sqrt{\frac{1}{\log \mu(B(t, \epsilon))}} d\epsilon \leq L\gamma_2(T, d). \quad (\text{A.10})$$

A probability measure μ on (T, d) such that the left-hand side of (A.10) is usefully small is called a majorizing measure. The (in)famous theory of majorizing measures used the infimum of the left-hand side of (A.10) over all choices of μ as a measure of the size of the metric space (T, d) .

Proof. Consider an admissible sequence (\mathcal{A}_n) with

$$\forall t \in T, \sum_{n \geq 0} 2^{n/2} \Delta(\mathcal{A}_n(t)) \leq 2\gamma_2(T, d).$$

Since $\text{card } \mathcal{A}_n \leq N_n$, there is a probability measure μ on T , supported by a countable set, and satisfying

$$\forall n \geq 1, \forall A \in \mathcal{A}_n, \mu(A) \geq \frac{1}{2^n N_n} \geq \frac{1}{N_n^2}$$

so that given $t \in T$

$$\begin{aligned} \epsilon > \Delta(\mathcal{A}_n(t), d) &\Rightarrow \mu(B(t, \epsilon)) \geq \frac{1}{N_n^2} \\ &\Rightarrow \sqrt{\log \frac{1}{\mu(B(t, \epsilon))}} \leq 2^{n/2+1}. \end{aligned} \quad (\text{A.11})$$

Now, since μ is a probability, $\mu(B(t, \epsilon)) = 1$ for $\epsilon > \Delta(T, d)$, and thus $\log(1/\mu(B(t, \epsilon))) = 0$. Thus

$$\begin{aligned} \int_0^\infty \sqrt{\log \frac{1}{\mu(B(t, \epsilon))}} d\epsilon &= \sum_{n \geq 1} \int_{\Delta(A_n(t))}^{\Delta(A_{n-1}(t))} \sqrt{\log \frac{1}{\mu(B(t, \epsilon))}} d\epsilon \\ &\leq \sum_{n \geq 1} 2^{n/2+1} \Delta(A_{n-1}(t)) \leq L\gamma_2(T, d) \end{aligned}$$

using (A.11). \square

Proposition A.4. *If μ is a probability measure on T , (supported by a countable set) and \mathcal{T} is a separated tree on T , then*

$$d(\mathcal{T}) \leq L \sup_{t \in T} \int_0^\infty \sqrt{\log \frac{1}{\mu(B(t, \epsilon))}} d\epsilon.$$

This completes the proof that the 4 “measures of the size of T ” considered in this appendix are indeed equivalent.

Proof. The basic observation is as follows. The sets

$$B(C, 4^{-s(A)-1}) = \{x \in T ; d(x, C) \leq 4^{-s(A)-1}\}$$

are disjoint as C varies over the children of A (as follows from (A.5)). So one of them has measure $\leq c(A)^{-1}$.

We then proceed in the following manner. We start with the largest element A_0 of \mathcal{T} . We then select a child A_1 of A_0 with $\mu(B(A_1, 4^{-s(A_0)-1})) \leq 1/c(A_0)$, and a child A_2 of A_1 with $\mu(B(A_2, 4^{-s(A_1)-1})) \leq 1/c(A_1)$, etc., and continue this construction as long as we can. It ends only when we reach a set of \mathcal{T} that has no child, and hence by (A.4) is reduced to a single point t . If $t \in A \in \mathcal{T}$, by construction we have

$$\mu(B(t, 4^{-s(A)-1})) \leq \frac{1}{c(A)}$$

so that

$$4^{-s(A)-2} \sqrt{\log c(A)} \leq \int_{4^{-s(A)-2}}^{4^{-s(A)-1}} \sqrt{\log \frac{1}{\mu(B(t, \epsilon))}} d\epsilon. \quad (\text{A.12})$$

By (A.6) the intervals $]4^{-s(A)-2}, 4^{-s(A)-1}[$ are disjoint for different sets A with $t \in A \in \mathcal{T}$, so summation of the inequalities (A.12) yields

$$\frac{1}{16} d(\mathcal{T}) \leq \sum_{t \in A \in \mathcal{F}} 4^{-s(A)-2} \sqrt{\log c(A)} \leq \int_0^\infty \sqrt{\log \frac{1}{\mu(B(t, \epsilon))}} d\epsilon.$$

\square

<http://www.springer.com/978-3-540-24518-6>

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Upper and Lower Bounds of Stochastic Processes

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2005, VIII, 222 p., Hardcover

ISBN: 978-3-540-24518-6