

Product Modularization Using Cuckoo Search Algorithm

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Abstract. Modularity plays an important role in managing complex systems by dividing them into a set of modules that are interdependent within and independent across the modules. In response to the changing market trend of having large varieties within small production processes, modular design has assumed significant roles in the product development process. The product to be designed is represented in the form of a Design Structure Matrix (DSM). The DSM contains a list of all product components and the corresponding dependency patterns among these components. The main objective of this paper is to support design of products under modularity through clustering products into a set of modules or clusters. In this research, Cuckoo Search optimization algorithm is used to find the optimal number of clusters and the optimal assignment of components to clusters. The objective is minimizing the total coordination cost. Results obtained show an improved performance compared to published studies.

Keywords: Design structure matrix · Cuckoo search · Modularity · Modular design · Clustering · Optimization

1 Introduction

System design involves clustering various components in a product such that the resulting modules are effective for the company. An ideal architecture is one that partitions the product into practical and useful modules. Some successfully designed modules can be easily updated on regular time cycles, some can be made in multiple levels to offer wide market variety, some can be easily removed as they stay, and some can be easily swapped to gain added functionality. The importance of effective product modularity is multiplied when identical modules are used in various different products [1].

Modularization in product design can help speed up the new product development process [2]. The product is represented in the form of a Design Structure Matrix (DSM) that contains a list of all product components and the corresponding information exchange and dependency patterns.

DSM, working as a product representation tool, provides a clear visualization of product design. The transformation of Component-DSM into proposed functional blocks of components is called clustering. For small problems' components, a Component-DSM may be sorted manually. For larger problems, this is not practical, and at some point, computer algorithms are absolutely necessary [3].

The aim of this paper is to develop a cuckoo search (CS) optimization algorithm to find: (1) the optimal number of clusters in a DSM; and (2) the optimal assignment of components to each cluster. The objective function is to minimize the total coordination cost. In this context, the DSM will work as a system analysis tool that provides a compact and clear representation of a complex system. It captures the interactions/interdependencies/interfaces between system elements. It also works as a project management tool which renders a project representation that allows for feedback and cyclic activity dependencies [4].

The rest of the paper is structured as follows. Section 2 provides a brief introduction on DSM. Section 3 reviews the literature and introduces the previous work this research builds on. Section 4 provides the problem definition. Section 5 presents the proposed algorithm. Section 6 presents and discusses the results obtained and, finally, Sect. 7 provides conclusion and ideas for future research.

2 Design Structure Matrix

The design structure matrix (DSM) is becoming a popular representation and analysis tool for system modeling. A DSM displays the relationships between components of a system or product in a compact visualization. Such a system can be, for example, product architecture or an engineering design process or a project.

The basic DSM is a simple square matrix, of size n , where n is the number of system elements. An example of a DSM is shown in Fig. 1. Element names are placed on the left hand side of the matrix as row headings and across the top row as column headings in the same order. If an element i depends on element j , then the matrix element ij (row _{i} , column _{j}) contains “1” or “x” otherwise the cell contains “0” or empty cell [5].

Once the DSM for a product is constructed, it can be analyzed for identifying modules, a process referred to as clustering. The goal of DSM clustering is to find a clustering arrangement where modules minimally interact with each other, while components within a module maximally interact with each other. As an example,

	A	B	C	D	E	F	G
A		x			x	x	
B				x			x
C		x		x			x
D		x	x		x		x
E				x		x	
F	x				x		
G		x	x	x			

(a)

	A	F	E	D	B	C	G
A		x	x		x		
F	x		x				
E		x		x			
D			x		x	x	x
B				x			x
C				x	x		x
G				x	x	x	

(b)

Fig. 1. Design structure matrix.

consider the DSM shown in Fig. 1(a). One can see from Fig. 1(b) that the DSM is rearranged by permuting rows and columns to contain most of the interactions within two separate modules: $\{A, F, E\}$ and $\{D, B, C, G\}$. However, three interactions are left out of any modules.

3 Related Work

The idea of maximizing interactions within modules and minimizing interactions between modules within a DSM was proposed by [6]. A stochastic clustering algorithm using this principle operating on a DSM was first introduced in [7], with subsequent improvements presented by [8]. The proposed algorithm can find clustering solutions to architecture and organization interaction problems modeled using DSM method. Gutierrez (1998) developed a mathematical model to minimize the coordination cost, and hence, find the optimal solution for a given number of clusters. A Simulated annealing algorithm was performed by [9] to find clustered DSM with cost minimization as an objective.

Yassine et al. (2007) used the design structure matrix (DSM) to visualize the product architecture, and to develop the basic building blocks required for the identification of product modules. The clustering method was based on the minimum description length (MDL) principle and a simple genetic algorithm (GA) [10].

Borjesson (2009) proposed a method for promoting better output from the clustering algorithm used in the conceptual module generation phase by adding convergence properties, a collective reference to data identified as option properties, geometrical information, flow heuristics, and module driver compatibility [11].

Van Beek et al. (2010) developed a modularization scheme based on the functional model of a system. The k-means clustering was adopted for DSM based modularization by defining a proper entity representation, a relation measure and an objective function [12]. A novel clustering method utilizing Neural Network algorithms and Design Structure Matrices (DSMs) was introduced by (Pandremenos and Chryssolouris 2012). The algorithm aimed to cluster components in DSM with predetermined number of clusters and clustering efficiency as an objective function [13].

Borjesson and Hölttä (2012) used IGTA (Idicula-Gutierrez-Thebeau Algorithm) for clustering Component-DSM as the basis for their work. They provided some improvement named IGTA-plus. IGTA-plus represented a significant improvement in the speed and quality of the solution obtained [3].

Borjesson and Sellgren (2013) presented an efficient and effective Genetic clustering algorithm, with the Minimum Description Length measure. To significantly reduce the time required for the algorithm to find good clusters, a knowledge aware heuristic element is included in the GA process. The efficiency and effectiveness of the algorithm is verified with four case studies [14].

Yang et al. (2014) provided a systematic clustering method for organizational DSM. The proposed clustering algorithm was able to evaluate the clustering structure based on the interaction strength [15].

Jung and Simpson (2014) introduced simple new metrics that can be used as modularity indices bounded between 0 and 1, and also utilized as the objective

functions to obtain the optimal DSM. The optimum DSM was the one with the maximized interactions within modules and the minimized interactions between modules [16].

Kim et al. (2015) provided a new approach for product design by integrating assembly and disassembly sequence structure planning [17].

The literature on DSM design shows that there are a few techniques available to cluster DSM for modularity. The main difference among these techniques is the clustering objective. Cost minimization is one of the first clustering objectives, where each DSM element is placed in an individual cluster, then, components are coordinated across modules. The objective is to minimize the cost of being inside and outside a cluster. A clustered DSM can be compared to a targeted DSM topology using another objective function called Minimal Description Length (MDL). MDL finds mismatching elements between the two topologies. The objective of clustering is to minimize MDL. The number of clusters is determined a priori based on the DSM structure. Clustering Efficiency (CE) index is another clustering objective that evaluates a weighed count of zero elements inside clusters and non-zero elements outside clusters with a predefined number of clusters.

4 Problem Definition

The problem presented in this work considers two decision variables: (1) the number of clusters to be formed and (2) the optimal assignment of elements to each cluster. The objective function is to minimize the total coordination cost. The total coordination cost of the DSM to be clustered is based on IntraClusterCost and ExtraClusterCost as shown in Eqs. 1 and 2,

$$\text{IntraClusterCost} = \sum_{i,k \in \text{Cluster } j} (\text{DSM}_{ik} + \text{DSM}_{ki}) * \text{Clustersize}(j)^{\text{powcc}}, \quad (1)$$

$$\begin{aligned} \text{ExtrClusterCost} = & \sum_{i,k \notin \text{cluster } j} (\text{DSM}_{ik} + \text{DSM}_{ki}) * \text{DSMSize}^{\text{powcc}}, \\ & j = 1 \dots \text{ncluster} \end{aligned} \quad (2)$$

where DSM_{ik} is the coupling between elements i and k , DSMSize is the number of elements (rows) in the matrix, powcc is the exponent used to penalize the size of clusters, and ncluster is the total number of clusters. clustersize is the number of elements in cluster j [18].

$$\text{Total coordination Cost} = \text{ExtraClusterCost} + \text{IntraClusterCost}$$

Subject to the constraint that each element is assigned only to one cluster, in other words, overlap between clusters is not allowed. Prohibiting overlap, or multi-cluster elements, is important for the following reasons: when allowing elements to be assigned in multiple clusters, the importance and usefulness of the clustering algorithm will be diminished or eliminated. If elements exist in more than one cluster, this forces

interactions between these clusters on multi levels. It is advantageous that elements placed in the same cluster are very similar [13].

Modularity affects both the profit and the sustainability of the product. A modular product contains modules that can be removed and replaced. The manufacturer can develop new modules instead of entirely new products. Therefore, customers buying upgraded modules only dispose of a portion of the product, thus reducing the total amount of waste. Hence, a customer upgrading a module does not have an entirely new product.

5 Proposed Algorithm

5.1 Cuckoo Search Algorithm and Pseudo Code

Yang and Deb (2009) proposed a new Meta heuristic algorithm called cuckoo search (CS). They tried to simulate the behavior of cuckoos to examine the solution space for optimization. The algorithm was inspired by the obligate interspecific brood parasitism of some cuckoo species that lay their eggs in the nests of host birds of other species. The aim is to escape the parental investment in raising their offspring. This strategy is also useful to minimize the risk of egg loss to other species, as the cuckoos can distribute their eggs amongst a number of different nests [19].

Of course, sometimes it happens that the host birds discover the alien eggs in their nests. In such case, the host bird takes different responsive actions varying from throwing such eggs away, to simply leaving the nest and building a new one elsewhere. On the other hand, the brood parasites have their own sophisticated characteristics to ensure that the host birds will care for the nestlings of their parasites. Examples of these characteristics are shorter egg incubation periods, rapid nestling growth, and egg coloration or pattern mimicking their hosts [20].

One major advantage of CS is its efficiency. The efficiency of the CS had been proven using many testing functions, for example, Michaelwicz function, Rosenbrock's function, etc. When comparing results with existing GA and PSO's, cuckoo search performed better [21]. Another major advantage of CS when compared to other metaheuristic algorithms, is its simplicity since it requires only two parameters. This feature reduces the effort of adjustment and fine tuning of parameter settings.

In cuckoo search, each egg can be regarded as a solution. In the initial process, each solution is generated randomly. When generating the i^{th} solution in $t + 1$ generation, denoted by X_i^{t+1} a levy flight is performed as shown in Eq. 3,

$$X_i^{t+1} = X_i^t + \alpha \oplus \text{Levy}(\lambda) \quad (3)$$

Where $\alpha > 0$ is a real number denoting the step size, which is related to the sizes of the problem of interest, and the $\oplus$ product denotes entry-wise multiplications. A Levy flight is a random walk where the step-lengths are distributed according to a heavy-tailed probability distribution as shown in Eq. 4. The random walk via Lévy flight is more efficient in exploring the search space as its step length is much longer in the long run.

$$\text{Levy} \sim u = t - \lambda, (1 < \lambda \leq 3) \quad (4)$$

The CS algorithm is based on three idealized rules [22].

- (1) Each cuckoo lays one egg at a time and dumps it in a randomly chosen nest.
- (2) The best nests with high quality eggs (solutions) will be carried over to the next generations.
- (3) The number of available host nests is fixed, and a host can discover an alien egg with a probability $pa \in [0, 1]$. In this case, the host bird can either throw the egg away or abandon the nest to build a completely new nest in a new location.

For simplicity, the third assumption can be approximated by a fraction pa of the n nests being replaced by new nests (with new random solutions at new locations). For a maximization problem, the quality or fitness of a solution could be proportional to the objective function. However, other more sophisticated expressions for the fitness function can also be defined.

Based on these three rules, the basic steps of the CS algorithm are summarized in the pseudo code in Fig. 2.

5.2 Solution Representation

The CS algorithm is used to solve the problem defined in Sect. 4. Solution representation of the problem is a vector of size equals to the number of elements in the DSM. Each cell in the vector takes an integer value between 1 and the number of elements, as

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Objective function  $f(x)$ ,  $x = (x_1, \dots, x_d)^T$ ;
Initial population of  $n$  host nests  $x_i$  ( $i = 1, 2, \dots, n$ );
while ( $t < \text{MaxGeneration}$ ) or (stop criterion);
  Get a cuckoo ( $i$ ) randomly using Levy flights;
  Evaluate its quality/fitness  $F_i$ ;
  Choose a nest among  $n$  ( $j$ ) randomly;
  if ( $F_i > F_j$ ),
    Replace  $j$  with the new solution;
  end
  Abandon a fraction ( $pa$ ) of worse nests
  and build new ones at new locations via Levy flights;
  Keep the best solutions (or nests with quality solutions);
  Rank the solutions and find the current best;
end while
Postprocess results and visualisation;

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Fig. 2. Pseudocode of CS [23].

1	2	2	2	3	3	1
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Fig. 3. Solution representation vector.

show in Fig. 3. The vector in Fig. 3 with size 7 represents a solution, where the DSM Contains 7 elements. Elements 1 and 7 belong to cluster 1, elements 2, 3, 4 belong to cluster 2, and elements 5, 6 belong to cluster 3. This solution representation forces the element to be a member of only one cluster.

Assume that we start with the maximum possible number of clusters, which equals to the number of elements in the DSM. The next step is to try to find the optimal number of clusters after deleting empty clusters. Such representation of the problem will not allow multi-clustering, which means each element will be assigned to only one cluster.

The problem is solved using Cuckoo search algorithm (CS). CS solves continuous types of variables. Since the problem in hand is categorized as a discrete variable problem, the solutions should be converted from continuous to discrete. This is done by the discretization of the continuous space by transforming the values into a limited number of possible states. There are several discretization methods available in the literature, for example: random key technique is used to transform from continuous space to discrete integer space. In order to decode the position, the nodes are visited in ascending order for each dimension [24]. Another discretization methods is the smallest position value (SPV) method. The technique maps the positions of the solution vector by placing the index of the lowest valued component as the first item on a permuted solution. The second lowest value component is the second item, and so on [25]. The nearest integer method is another technique, to transform continuous variables to integer variables. In the nearest integer method, a real value is converted to the nearest integer (NI) by rounding or truncating up or down [26].

Considering the above mentioned methods, SPV, and random key methods, are not suitable for the problem presented in this work. This is because integer value(s) need to be repeated, while these methods result in unique values. Therefore, the suitable method for the problem in hand is the nearest integer method since it allows the repetition of values by truncating to the higher or lower value.

To cluster a DSM into modules using CS, we start with a set of nests; each nest is a vector of length that equals to the number of elements in the DSM. This vector contains random numbers following uniform distribution in the range from lower and upper limits. These random numbers are converted to integer values using the nearest integer method. Each one of these integer numbers represent a solution that could be sent for the evaluation function. The evaluation function returns the total coordination cost.

5.3 Solution Evaluation

The total coordination cost of the DSM to be clustered is based on IntraClusterCost and ExtraClusterCost as explained in Sect. 4. Regarding intracluster cost, if interaction DSM_{ik} belongs to cluster j , then calculate intra cluster cost. On the other hand, if

interaction DSM_{ik} does not belong to cluster j , calculate the extra cluster cost. The first step in calculating the total coordination cost is to start with the total number of interactions in the DSM multiplied by the size of the DSM raised to the power $powcc$. This is the highest value of total coordination. This value will be minimized in subsequent steps of the algorithm after forming clusters. After completion of the evaluation step, select the best solution and go to the next best solution using Levy flight, carrying the best nests with high quality eggs (solutions) over to the next generations. Continue till the stopping condition is reached.

6 Experimental Results and Analysis

In this section we examine the CS algorithm on a number of instances. The CS algorithm has two parameters namely, the (pa) value. The (pa) value represents the probability of discovering an alien egg, which corresponds to getting rid of solution and the number of nests corresponds to the population of solutions. There are no general recommendations in the literature on the range(s) of values for these two parameters. Therefore, we examine the solution quality of the test instances over a range of possible values for the two parameters. We chose the different values of (pa) in the range [0.1, 0.9], with an increment of 0.1. We also examine the solution quality with different number of nests, namely, 25, 50, 75, and 100. We noticed that the change in the number of nests did not have a significant impact on the objective function value in all tested cases. The maximum number of iterations is 3000 and it is used as the stopping condition. The algorithm is coded using Matlab, and is run on a computer with Intel(R) Core(TM) i3 CPU, 2.27 GHz PC. We use three test instances, 2 are categorized as small size instances with 7 and 9 elements each, available in [9, 10], respectively. The third instance is larger in size and consists of 61 elements, and is available in [9].

The first small size instance has a DSM that contains 7 elements as shown in Fig. 4. The DSM starts initially with a total coordination cost of 68. This cost is based on assigning each element in it is own cluster. No clusters are formed yet.

After applying the CS clustering algorithm, the clustered DSM is as shown in Fig. 5. The Total coordination cost is reduced to 48. This problem was solved in [9] and a total coordination cost of 53 was obtained. Hence, the proposed CS is able to obtain better results. The minimum number of clusters using the proposed CS algorithm is 2. Figure 6 shows the total cost as it changes with every iteration. The best solution is obtained in iteration number 575. The CPU run time ranges from 0.07 s to 0.66 s for 100 to 3000 iterations, respectively.

It is noticed that, in the clustered DSM two clusters are formed and the majority of the interactions are included in clusters. This indicates that similar elements are grouped in the same cluster. In this case, IntraClusterCost is larger than the ExtraClusterCost which improves the objective function value. Only 3 interactions are placed outside clusters (number of 1's).

Figure 7 shows the objective function values corresponding to different (pa) values. It is noticed from Fig. 7 that the objective function value remains 48 till the value of (pa) reaches 0.5. When the value of (pa) exceeds 0.5, the objective function value

	1	2	3	4	5	6	7
1		1	0	0	1	1	0
2	0		0	1	0	0	1
3	0	1		1	0	0	1
4	0	1	1		1	0	1
5	0	0	0	1		1	0
6	1	0	0	0	1		0
7	0	1	1	1	0	0	

Fig. 4. Original DSM.

	1	5	6	2	3	4	7
1		1	1	1	0	0	0
5	0		1	0	0	1	0
6	1	1		0	0	0	0
2	0	0	0		0	1	1
3	0	0	0	1		1	1
4	0	1	0	1	1		1
7	0	0	0	1	1	1	

Fig. 5. Clustered DSM.

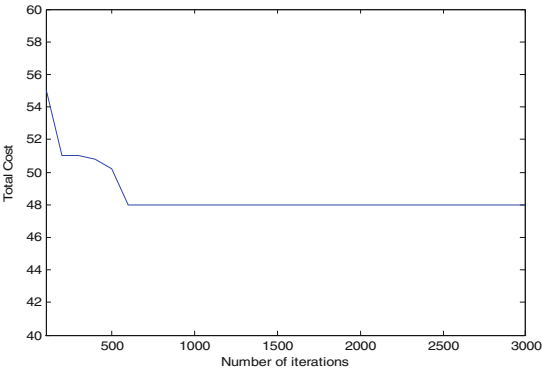


Fig. 6. Cost history for CS algorithm-best solution.

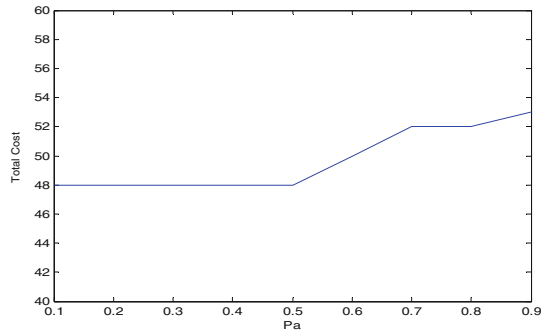


Fig. 7. The relation between P_a and total cost.

deteriorates. This is due to the fact that high values of (p_a) tend to get rid of solutions without trying to improve them locally. Therefore, the value of (p_a) in the range [0.1, 0.5] for this test instance, achieves the required balance between exploration and exploitation.

We examined the developed algorithm on another problem presented in [10]. The DSM of the problem has 9 elements as shown in Fig. 8.

Figure 9 shows the clustered DSM after using CS algorithm. The total coordination cost after clustering with CS is 41.8. The corresponding number of clusters is 4. The resulting DSM clustered using our proposed CS algorithm is the same as the one obtained in [10]. Figure 10 shows the total cost as it changes with every iteration. The best solution is obtained in iteration number 880. The CPU run time ranges from 1.07 s to 13.52 s for 100 to 3000 iterations, respectively.

We notice from Fig. 9 that, in the clustered DSM four clusters are formed, cluster 1 with the most similar 3 elements, cluster 2 with the most similar 4 elements, cluster 3 with 1 element and cluster 4 with 1 element. All interactions are included in clusters. In this case, there are no extra costs because no 1's are outside clusters.

	A	B	C	D	E	F	G	H	I
A	1	0	0	0	1	0	1	0	0
B	0	1	1	0	0	1	0	1	0
C	0	1	1	0	0	1	0	1	0
D	0	0	0	1	0	0	0	0	0
E	1	0	0	0	1	0	1	0	0
F	0	1	1	0	0	1	0	1	0
G	1	0	0	0	1	0	1	0	0
H	0	1	1	0	0	1	0	1	0
I	0	0	0	0	0	0	0	0	1

Fig. 8. Original DSM.

	A	E	G	B	C	F	H	D	I
A		1	1						
E	1		1						
G	1	1							
B					1	1	1		
C				1		1	1		
F				1	1		1		
H				1	1	1			
D									
I									

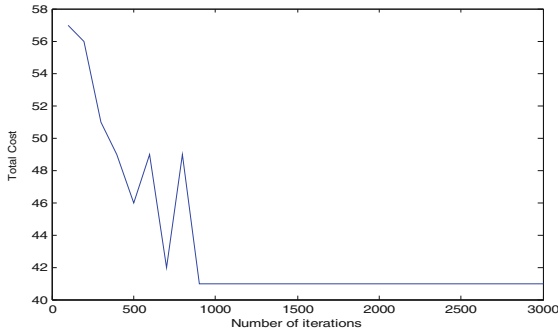
Fig. 9. Clustered DSM.**Fig. 10.** Cost history for CS algorithm-best solution.

Figure 11 shows the objective function values corresponding to different (pa) values. It is noticed from Fig. 11 that the objective function value remains 41.8 till the value of (pa) reaches 0.6. When the value of (pa) exceeds 0.6, the objective function value deteriorates. The impact of the value of (pa) on solution quality is similar to that in the first test instance.

To further evaluate the proposed CS algorithm we apply it on a large size problem, available in [9]. The DSM contains 61 elements and represents an elevator example. The total coordination cost obtained using the CS algorithm is 4133.25, with a total number of 17 clusters. The total coordination cost obtained in [9] is 4433. Accordingly, our proposed CS algorithm is able to obtain superior results when compared to the results obtained by [9]. Cluster assignments of the elevator example using the CS algorithm are shown in Table 1. The best solution is obtained in iteration number 921. The CPU run time is 450.7 s after 3000 iterations. Figure 12 shows the total cost as it changes with every iteration.

Table 1 shows that 17 clusters are formed, each cluster contains the most similar elements which minimizes the total coordination cost. Figure 13 shows the impact of

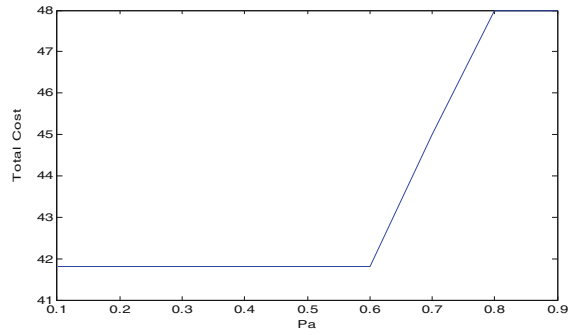


Fig. 11. The relation between P_a and total cost.

changing the values of (p_a) on the objective function value. Similar to the previous test instances, it is noticed that the solution deteriorates as the value of (p_a) exceeds a certain value. For this test instance the (p_a) value after which the solution starts to deteriorate is 0.4.

Table 1. Results obtained using CS algorithm for the elevator example.

Cluster number	Elements that cluster contains
1	1, 3, 5, 11, 15, 17, 18, 20, 22, 28, 34, 35, 37, 39, 40, 41, 43, 47, 48, 50, 59, 60, 61
2	2, 8, 12, 16, 19, 21, 26, 27, 32, 33, 44, 46, 49, 54
3	4
4	6, 9, 13
5	7
6	10, 14, 25, 55
7	20
8	22, 53
9	23, 31
10	24
11	30
12	36
13	51
14	52
15	56
16	57
17	58

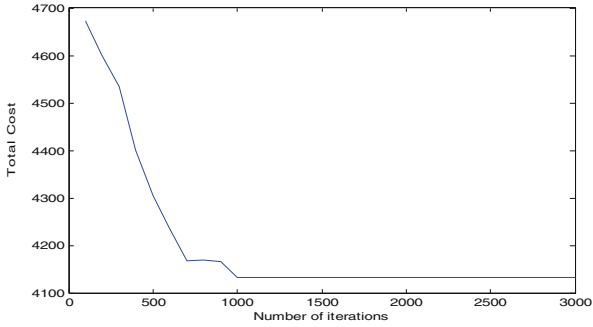


Fig. 12. Cost history for CS algorithm–best solution.

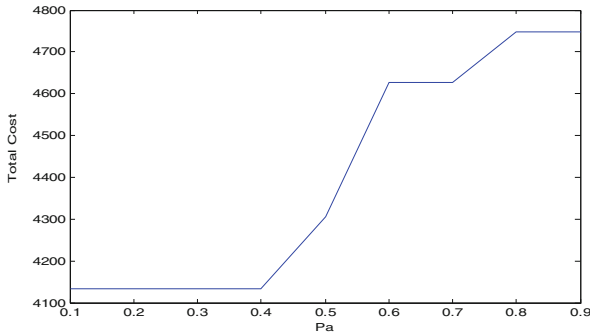


Fig. 13. The relation between Pa and total cost.

7 Conclusion and Future Work

The importance of modular product design has been recognized due to many reasons, for example, its ability to achieve economies of scale, providing high degree of flexibility during production, and reducing lead times. Clustering of a design structure matrix (DSM) is one of the methods that are widely used to design a product under modularity. In this work, we aimed at designing modular products through representation and clustering of their DSM. The clustering process required solving for 2 decision variables: the number of clusters and the elements assigned to each cluster. The objective function was minimizing the total coordination cost, under the constraint that each element was to be assigned to one cluster, without allowing clusters' overlap. Heuristics and metaheuristic techniques are typically used to cluster DSMs for modularity. Cuckoo search (CS) is a metaheuristic algorithm that has proven efficiency in terms of solution quality, speed, and simplicity, when compared to other metaheuristics algorithms available in the literature. In this work, we adopted CS algorithm to perform DSM clustering. To test its performance, we applied CS on a number of instances available in the literature. The proposed CS algorithm was able to obtain better or similar results for all the test instances. Future work includes restricting the

number of elements belonging to each cluster, considering inventory and supply chain decisions in the modular design process, and including the number of clusters in the objective function.

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