

The Classification of All the Subvarieties of DNMG

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Abstract. MTL is the logic of all left-continuous t -norms and their residua. The equivalent algebraic semantics of MTL is constituted by the variety of MTL-algebras, MTL . The variety WNM of *weak nilpotent minimum* algebras is a major subvariety of MTL , containing several subvarieties of MTL which have been subjects of study in the literature, such as Gödel algebras, Nilpotent Minimum algebras, Drastic Product and Revised Drastic Product algebras, NMG -algebras, as well as Boolean algebras. In this paper we introduce and axiomatise DNMG , a proper subvariety of WNM which contains all the aforementioned varieties. We show that DNMG is singly generated by a standard algebra. Further, we determine the structure of the lattice of subvarieties of DNMG , and we provide the axiomatisation of every subvariety.

Keywords: WNM -algebras · DNMG -algebras · NM -algebras · Gödel-algebras · DP -algebras · Axiomatisations of subvarieties · Single chain completeness

1 Introduction

Nilpotent minimum t -norm $*_{\mathsf{NM}}$ [14] was one of the first examples of a left-continuous but not continuous t -norm. The logic related to $*_{\mathsf{NM}}$, NM , was introduced by Esteva and Godo in [12]. In the same paper they presented a generalisation of NM , the logic of Weak Nilpotent Minimum, WNM , and the related algebraic semantics, the variety of WNM -algebras WNM . WNM is an extension of MTL , the logic of all left continuous t -norms and their residua [12, 20]. Several extensions of WNM have been extensively studied in the literature. In particular, Gödel logic G , Drastic Product DP ([4], firstly introduced as $\mathsf{S}_3\mathsf{MTL}$ in [23]), Revised Drastic product RDP ([9, 25], based on the t -norm introduced in [19]), NMG [26], NM [12], and classical Boolean logic B . During the years a number of topics concerning WNM and its algebraic semantics has been investigated: the papers [5, 13, 16, 21, 24] are only few examples. WNM has been extensively studied in [23], where the problem of axiomatising its extensions has been partially

solved. The task of characterising and axiomatising the lattice of extensions of a given extension of WNM has been accomplished in some cases. Gispert [17], solved the case for NM . The lattice of extensions of G is well known. The extensions of EMTL , which is the largest common fragment of G and DP , have been axiomatised in [6].

In this article we introduce the variety $\mathbb{D}\text{NMG}$, the algebraic semantics of DNMG , which is an extension of WNM which is a particularly tame single-chain complete common fragment of G , NM and DP (and hence, of all the aforementioned extensions of WNM). We prove standard completeness for DNMG . Generalising Gispert's result, we describe the structure of the lattice of subvarieties of $\mathbb{D}\text{NMG}$, showing that each one of them is generated by finitely many chains. We further provide a uniform way to axiomatise each one of these subvarieties.

2 Preliminaries

We assume that the reader is acquainted with many-valued logics in Hájek's sense, and with their algebraic semantics. We refer to [11, 18] for any unexplained notion. We recall that MTL is the logic, on the language $\{\&, \wedge, \vee, \rightarrow, \neg, \perp, \top\}$, of all left-continuous t -norms and their residua, and that its associated algebraic semantics in the sense of Blok and Pigozzi [7] is the variety MTL of *MTL-algebras*, that is, prelinear, commutative, bounded, integral, residuated lattices [11]. In an MTL -algebra $\mathcal{A} = (A, *, \Rightarrow, \sqcap, \sqcup, \sim, 0, 1)$ the connectives $\&, \rightarrow, \wedge, \vee, \neg, \perp, \top$ are interpreted, respectively, by $*, \Rightarrow, \sqcap, \sqcup, \sim, 0, 1$. Totally ordered MTL -algebras are called MTL -chains. In every chain $\sqcap = \min$ and $\sqcup = \max$. An MTL -algebra is called standard whenever its lattice reduct is $([0, 1], \leq)$, with the usual order.

Given an MTL -algebra \mathcal{A} , with $\mathbf{V}(\mathcal{A})$ we mean the variety generated by \mathcal{A} , which is said to be *generic* for $\mathbf{V}(\mathcal{A})$. A logic L is the extension of MTL via a set of axioms $\{\varphi_i\}_{i \in I}$ if and only if \mathbb{L} is the subvariety of MTL -algebras satisfying $\{\bar{\varphi}_i = 1\}_{i \in I}$, where $\bar{\varphi}_i$ is obtained from φ_i by replacing the connectives with the corresponding operations, and every propositional variable in φ with an individual variable. With $\mathcal{A} \models \bar{\varphi} = 1$ we mean that \mathcal{A} satisfies $\bar{\varphi} = 1$.

The logic WNM [12] is axiomatised as MTL plus:

$$\neg(\varphi \& \psi) \vee ((\varphi \wedge \psi) \rightarrow (\varphi \& \psi)). \quad (\text{wnm})$$

The logics G , DP , EMTL , RDP , NM , NMG [1, 4, 6, 12, 25, 26] are axiomatised as WNM plus, respectively:

$$\varphi \rightarrow (\varphi \& \varphi). \quad (\text{id})$$

$$\varphi \vee \neg(\varphi \& \varphi). \quad (\text{dp})$$

$$(\varphi \rightarrow (\varphi \& \varphi)) \vee (\psi \vee \neg(\psi \& \psi)). \quad (\text{emtl})$$

$$(\varphi \rightarrow \neg\varphi) \vee \neg\neg\varphi. \quad (\text{rdp})$$

$$\neg\neg\varphi \rightarrow \varphi. \quad (\text{inv})$$

$$\neg\neg\varphi \vee (\neg\neg\varphi \rightarrow \varphi). \quad (\text{nmg})$$

NM^- [17] is axiomatised as NM plus $\neg((\neg(\varphi^2))^2) \leftrightarrow (\neg((\neg\varphi)^2))^2$, where φ^2 stands for $\varphi \& \varphi$. \mathbf{B} is classical logic, axiomatised as MTL plus $\varphi \vee \neg\varphi$.

As shown in [17], the operations $*$, \Rightarrow of a WNM -chain \mathcal{A} , are:

$$x * y = \begin{cases} 0 & \text{if } x \leq \sim y \\ \min(x, y) & \text{Otherwise.} \end{cases} \quad x \Rightarrow y = \begin{cases} 1 & \text{if } x \leq y \\ \max(\sim x, y) & \text{Otherwise.} \end{cases} \quad (1)$$

Where \sim , the negation of \mathcal{A} , is a (*generalised*) *weak negation*, that is a map $\sim : A \rightarrow A$ such that $\sim 1 = 0$, $\sim \sim a \geq a$, and if $a \leq b$, then $\sim a \geq \sim b$. Each weak negation is the negation of a uniquely determined WNM -chain. WNM is locally finite, i.e. for every WNM -algebra each one of its finitely generated subalgebras is finite. For each integer $n \geq 2$, with \mathbb{G}_n , \mathbb{DP}_n , NM_{2n}^- we will denote the variety generated, respectively, by the Gödel chain with n elements, the DP -chain with n elements, and the NM^- -chain with $2n$ elements.

3 DNMG-algebras

DNMG is the variety of WNM -algebras satisfying the following identity.

$$\sim \sim x \sqcup (\sim \sim x \Rightarrow x) \sqcup (\sim \sim x \Leftrightarrow \sim x) = 1. \quad (\text{DNMG})$$

Since each variety of MTL -algebras is generated by its chains, we immediately have that two subvarieties of MTL -algebras coincide iff they have the same class of chains. We are then going to analyse the structure of DNMG -chains.

Definition 1. Let \mathcal{A} be a WNM -chain. Let us define the following sets.

- $A^+ = \{a \in A : a > \sim a\}$, and $A^- = \{a \in A : a < \sim a\}$.
- $S(A) = \{a \in A \mid \sim \sim a = 1, a \neq 1\}$.
- $F(A) = \{a \in A \mid \sim a = \sim \sim a\}$.
- $I^-(A) = \{a \in A \mid \sim \sim a = a, 0 < a < \sim a\}$.
- $I^+(A) = \{a \in A \mid \sim \sim a = a, 1 > a > \sim a\}$.
- $I(A) = I^-(A) \cup I^+(A)$.

Clearly, $I^-(A) \cap I^+(A) = \emptyset$. Further, \sim is a bijection of $I(A)^+$ onto $I(A)^-$. Notice that $S(A)$ is disjoint from $I(A)$ and $F(A)$. Given $B, C \subseteq A$, we write $B \prec C$ whenever $b <_A c$, for every $b \in B$ and $c \in C$. The following is immediate.

Proposition 1. For every WNM -chain \mathcal{A} it holds that $I^-(A) \prec F(A) \prec I^+(A) \prec S(A)$. In particular, $I(A)^- \cup F(A) \cup \{0\} = A \setminus A^+$, and $I(A)^+ \cup S(A) \cup \{1\} = A^+$.

By Proposition 1, the sets $S(A), F(A), I^-(A), I^+(A)$ are pairwise disjoint. For any subset $S \subseteq A$, let $\langle S \rangle$ be the subalgebra of \mathcal{A} generated by S .

Proposition 2. *Let \mathcal{A} be a WNM-chain and pick $S \subseteq A$.*

1. *If $S \subseteq I(A)$, then $\langle S \rangle = S \cup \{\sim a \mid a \in S\} \cup \{0, 1\}$ is an NM^- algebra.*
2. *If $S \subseteq F(A)$, then $\langle S \rangle = S \cup \{\sim a \mid a \in S\} \cup \{0, 1\}$ is a DP-algebra.*
3. *If $S \subseteq S(A)$ then $\langle S \rangle = S \cup \{0, 1\}$ is a Gödel algebra.*

Proof. Using (1) an easy check shows that the sets $(S \cap I(A)) \cup \{\sim a \mid a \in S \cap I(A)\} \cup \{0, 1\}$, $(S \cap F(A)) \cup \{\sim a \mid a \in S \cap F(A)\} \cup \{0, 1\}$ and $(S \cap S(A)) \cup \{0, 1\}$ are subuniverses of A . The rest follows by Definition 1. \square

Theorem 1. *A WNM-chain \mathcal{A} is a DNMG-chain iff $A = S(A) \cup F(A) \cup I(A) \cup \{0, 1\}$.*

Proof. Let \mathcal{A} be a WNM-chain such that $A = S(A) \cup F(A) \cup I(A) \cup \{0, 1\}$. Then, each element $a \in S(A)$ satisfies $\sim \sim a = 1$, each element $a \in F(A)$ satisfies $\sim \sim a = \sim a$, and each element $a \in I(A)$ satisfies $\sim \sim a = a$. The elements 0 and 1 both satisfy $\sim \sim a = a$. Whence \mathcal{A} satisfies the identity (DNMG). Conversely, let \mathcal{A} be a WNM-chain such that there is $a \in A \setminus (S(A) \cup F(A) \cup I(A) \cup \{0, 1\})$. Then $\sim \sim a \neq 1$, $\sim \sim a \neq \sim a$ and $\sim \sim a \neq a$. Whence \mathcal{A} is not a DNMG-chain. \square

Lemma 1. *Let $\mathbb{V} \in \{\mathbb{DP}, NM^-, \mathbb{G}\}$. For each subvariety \mathbb{W} of \mathbb{V} , a chain \mathcal{A} is generic for \mathbb{W} iff each finite chain in \mathbb{W} embeds into \mathcal{A} . Moreover, $\mathbf{V}(\mathcal{A}) = \mathbb{V}$ iff \mathcal{A} is infinite.*

Proof. Let $\mathbb{V} \in \{\mathbb{DP}, NM^-, \mathbb{G}\}$, and pick two chains $\mathcal{A}, \mathcal{B} \in \mathbb{V}$. By the results of [6, 17] we have two consequences. First, $\mathbf{V}(\mathcal{A}) = \mathbb{V}$ iff \mathcal{A} is infinite. Next, if $|\mathcal{A}| < |\mathcal{B}|$ then \mathcal{A} embeds into \mathcal{B} . The claim follows immediately. \square

Definition 2. *Let $C \subseteq \{F, I, S\}$, and let \mathcal{A} be a WNM-chain. We say that \mathcal{A} is C -semigeneric iff for any $X \in C$ and for any finite chain $\mathcal{B} \in \mathbb{WNM}$, it holds that $\langle X(\mathcal{B}) \rangle$ embeds into $\langle X(\mathcal{A}) \rangle$.*

In the following, for any $X \in \{F, I, S\}$, by $L(X)$ -chain we mean: DP-chain if $X = F$, NM^- -chain if $X = I$, and G-chain if $X = S$.

Lemma 2. *For each $X \in \{F, I, S\}$, a WNM-chain \mathcal{A} is X -semigeneric iff $\langle X(\mathcal{A}) \rangle$ is generic for the variety generated by $L(X)$ -chains.*

Proof. By Lemma 1 and Proposition 2. \square

Definition 3. *For each $n > 0$, we let $e_n(F)$, $e_n(I)$ and $e_n(S)$ denote the following terms*

$$\begin{aligned}
 e_n(F) &= \bigsqcup_{i=1}^n ((\sim \sim x_i) \sqcup ((\sim \sim x_i \Rightarrow x_i) \sqcap (\sim((\sim(x_i^2))^2) \Leftrightarrow (\sim((\sim x_i)^2))))^2), \\
 e_n(I) &= \bigsqcup_{i=1}^n ((\sim \sim x_i) \sqcup (\sim x_i) \sqcup (\sim \sim x_i \Leftrightarrow \sim x_i)), \\
 e_n(S) &= \bigsqcup_{i=1}^n ((\sim \sim x_i \Leftrightarrow \sim x_i) \sqcup (\sim \sim x_i \Rightarrow x_i)).
 \end{aligned}$$

Lemma 3. *Let \mathcal{A} be a DNMG-chain, $X \in \{F, I, S\}$ and, for each $n > 0$, let $(a_1, a_2, \dots, a_n) \in A^n$. Then, $(e_n(X))(a_1, \dots, a_n) < 1$ iff $(a_1, \dots, a_n) \in X(A)$.*

Proof. By Definition 1, it is sufficient to note that the i th disjunct of $e_n(X)$ evaluates to 1 iff $a_i \in A \setminus X(A)$. \square

Lemma 4. *Let \mathcal{A} be a non-trivial DNMG-chain, and $t(x_1, \dots, x_n) = 1$ be an equation. Then, for every $X \in \{F, I, S\}$, the equation $t = 1$ holds in $\langle X(A) \rangle$ iff $t \sqcup e_n(X) = 1$ holds in \mathcal{A} .*

Proof. If $X(A) = \emptyset$, then $\langle X(A) \rangle$ is isomorphic with the two-element Boolean algebra $\{0, 1\}$, and by Lemma 3 $\mathcal{A} \models e_n(X) = 1$. The claim follows immediately. We then assume $X(A) \neq \emptyset$: by Lemma 3 we have that $\langle X(A) \rangle \models e_n(X) = 1$. Assume first that $\mathcal{A} \models t \sqcup e_n(X) = 1$: we must have $\langle X(A) \rangle \models t = 1$. Conversely, suppose $\mathcal{A} \not\models t \sqcup e_n(X) = 1$: by Lemma 3 we have that for some $a_1, \dots, a_n \in \langle X(A) \rangle$, $e_n(X)(a_1, \dots, a_n) < 1$ and $t(a_1, \dots, a_n) < 1$. We conclude that $\mathcal{A} \not\models t = 1$. \square

Lemma 5. *A DNMG-chain \mathcal{A} embeds into a DNMG-chain \mathcal{B} iff $\langle X(A) \rangle$ embeds into $\langle X(B) \rangle$ for each $X \in \{F, I, S\}$.*

Proof. One direction is trivial. Assume then that there is $X \in \{F, I, S\}$ such that $\langle X(A) \rangle$ does not embed into $\langle X(B) \rangle$, and assume further, by contradiction, that $f: A \rightarrow B$ is an embedding of \mathcal{A} into \mathcal{B} . Then there is $a \in \langle X(A) \rangle$ such that $f(a) \notin \langle X(B) \rangle$, for otherwise f would embed $\langle X(A) \rangle$ into $\langle X(B) \rangle$. Assume $X = F$ and $f(a) \in B \setminus \langle F(B) \rangle$. Then $\sim \sim a = \sim a$, while $\sim \sim f(a) \neq \sim f(a)$, contradicting the fact that f is an homomorphism. The other cases $X = I$ or $X = S$, are dealt with analogously. \square

Lemma 6. *Let \mathbb{V} be a subvariety of DNMG. Then a DNMG-chain $\mathcal{A} \in \mathbb{V}$ is generic for \mathbb{V} iff each finite chain \mathcal{B} in \mathbb{V} embeds into \mathcal{A} .*

Proof. Assume that each finite chain \mathcal{B} in \mathbb{V} embeds into \mathcal{A} . Since \mathbb{V} has the finite model property, being locally finite, then it is generated by the class of its finite chains. This implies that $\mathbf{V}(\mathcal{A}) = \mathbb{V}$. On the other hand, assume that \mathcal{B} does not embed into \mathcal{A} . Then, by Lemma 5, there is $X \in \{F, I, S\}$ such that $\langle X(B) \rangle$ does not embed into $\langle X(A) \rangle$. By Lemma 1, $\langle X(B) \rangle$ is an $L(X)$ -chain which does not belong to the variety generated by the $L(X)$ -chain $\langle X(A) \rangle$. Whence there is an equation $t(x_1, \dots, x_n) = 1$ holding in $\langle X(A) \rangle$ and failing in $\langle X(B) \rangle$. Then, by Lemma 4, the equation $t \sqcup e_n(X) = 1$ holds in \mathcal{A} and fails in \mathcal{B} , proving that \mathcal{A} is not generic for \mathbb{V} . \square

Lemma 7. *Let $\mathbb{V} \subseteq \text{DNMG}$ be a single chain generated variety. Then there is a chain $\mathcal{B} \in \mathbb{V}$ such that every countable chain in \mathbb{V} embeds into it, and $\mathbf{V}(\mathcal{B}) = \mathbb{V}$.*

Proof. Immediate by Lemma 6, [10, Theorem 3.8] and [22, Theorems 3 and 5].

Lemma 8. *A DNMG-chain is $\{F, I, S\}$ -semigeneric iff it is generic for DNMG.*

Proof. Immediate, from Lemmas 2, 5 and 6. \square

Theorem 2. *A chain \mathcal{A} is generic for DNMG iff the sets $F(\mathcal{A})$, $I(\mathcal{A})$ and $S(\mathcal{A})$ are infinite.*

Proof. By Lemmas 1 and 8. □

We are ready to prove standard completeness for the logic DNMG whose associated algebraic semantics is given by DNMG.

Definition 4. *Let $*$: $[0, 1]^2 \rightarrow [0, 1]$ be defined by:*

$$x * y = \begin{cases} 0 & \text{if } x + y \leq 3/4 \text{ or } \max\{x, y\} \leq 1/2 \\ \min\{x, y\} & \text{otherwise.} \end{cases}$$

It is immediate to check that $*$ is a left-continuous t -norm. Moreover $*$ is the monoidal operation of the standard WNM-chain $[0, 1]_*$ determined by the following negation: $\sim 0 = 1$, $\sim x = 3/4 - x$ for all $x \in (0, 1/4) \cup (1/2, 3/4]$, $\sim x = 1/2$ for all $x \in [1/4, 1/2]$, $\sim x = 0$ for all $x \in [3/4, 1]$.

Lemma 9. *The WNM-chain $[0, 1]_*$ determined by the t -norm in Definition 4 is a DNMG-chain.*

Proof. We just check that any element $a \in [0, 1/4) \cup (1/2, 3/4]$ is such that $\sim \sim a = a$, while any element $a \in [1/4, 1/2]$ is such that $\sim \sim a = \sim a = 1/2$, and finally any element $a \in [3/4, 1]$ is such that $\sim \sim a = 1$. □

Theorem 3. *The logic DNMG is standard complete, since the variety DNMG is generated by $[0, 1]_*$.*

Proof. By Lemma 9, $[0, 1]_* \in \text{DNMG}$. Now, $F([0, 1]_*) = [1/4, 1/2]$, $I([0, 1]_*) = (0, 1/4) \cup (1/2, 3/4)$, and $S([0, 1]_*) = [3/4, 1]$. By Lemma 8 and Theorem 2, the standard chain $[0, 1]_*$ is generic for DNMG. □

4 The Lattice of Subvarieties of DNMG

Let ω denote the ordinal $\{0, 1, 2, \dots\}$ of the natural numbers, and let $\omega + 1$ be the ordinal $\omega \cup \{\omega\}$. For any integer $n > 0$ and any sequence $\kappa_1, \kappa_2, \dots, \kappa_n$ of ordinals, the *direct product* $\kappa_1 \times \kappa_2 \times \dots \times \kappa_n$ is the poset obtained equipping the cartesian product with the pointwise order: $(a_1, \dots, a_n) \leq (b_1, \dots, b_n)$ iff $a_i \leq b_i$ for all $i \in \{1, 2, \dots, n\}$. We write $\kappa^{(n)}$ to mean the n th *direct power* of the ordinal κ , that is, the direct product of n copies of κ .¹ A subset $S \subseteq \kappa_1 \times \dots \times \kappa_n$ is an *antichain* if for each $a, b \in S$, neither $a \leq b$ nor $b \leq a$. The set of all antichains of a poset P is denoted $\mathcal{AC}(P)$.

Lemma 10. *For $n > 0$, every antichain of $P = \kappa_1 \times \kappa_2 \times \dots \times \kappa_n$ is finite.*

¹ We use this notation to distinguish direct powers from ordinal exponentiation.

Proof. By [8, Exercise 2.3.4] a poset is a well-quasi-order (wqo) iff (1) it has no infinite strictly decreasing sequences, i.e. $x_0 > x_1 > \dots$, and (2) it has no infinite antichains. As every ordinal κ_i is well-ordered, then it satisfies (1) and (2): whence it is a wqo. By [8, Lemma 2.3.9] P is also a wqo, whence each one of its antichains is finite. \square

Corollary 1. *For each integer $n > 0$, the set $\mathcal{AC}((\omega + 1)^{(n)})$ has cardinality \aleph_0 .*

Proof. By Lemma 10, every antichain of $(\omega + 1)^{(n)}$ is a finite set of n -tuples of $\omega \cup \{\omega\}$. Hence it can be coded in the natural numbers, i.e. there is an injective map $\mathcal{AC}((\omega + 1)^{(n)}) \rightarrow \omega$. Whence, $|\mathcal{AC}((\omega + 1)^{(n)})| \leq \aleph_0$. To conclude, note that $\mathcal{AC}((\omega + 1)^{(1)}) = \omega + 1$, whence, $|\mathcal{AC}((\omega + 1)^{(n)})| \geq |\mathcal{AC}((\omega + 1)^{(1)})| = \aleph_0$. \square

The set $\mathcal{AC}((\omega + 1)^{(n)})$ can be equipped with a lattice structure by putting $X \leq Y$ if for each n -tuple $x \in X$ there is $y \in Y$ such that $x \leq y$, for all $X, Y \in \mathcal{AC}((\omega + 1)^{(n)})$. In this section we shall prove that the lattice of subvarieties of DNMG is isomorphic with $\mathcal{AC}((\omega + 1)^{(3)})$.

Definition 5. *Let \mathcal{A} be a DNMG-chain. Then the triplet $T(\mathcal{A})$ associated with \mathcal{A} is an element $(a, b, c) \in (\omega + 1)^{(3)}$ defined as follows.*

1. *If $S(\mathcal{A})$ is infinite then $a = \omega$, otherwise $a = |S(\mathcal{A})|$.*
2. *If $I^-(\mathcal{A})$ is infinite then $b = \omega$, otherwise $b = |I^-(\mathcal{A})|$.*
3. *If $F(\mathcal{A})$ is infinite then $c = \omega$, otherwise $c = |F(\mathcal{A})|$.*

Lemma 11. *Let \mathcal{A} and \mathcal{B} be two DNMG-chains with \mathcal{A} of finite cardinality. Then \mathcal{A} embeds into \mathcal{B} iff $T(\mathcal{A}) \leq T(\mathcal{B})$ in the pointwise order.*

Proof. \mathcal{A} embeds into \mathcal{B} iff, by Lemma 5, $\langle S(\mathcal{A}) \rangle$ embeds into $\langle S(\mathcal{B}) \rangle$, $\langle I(\mathcal{A}) \rangle$ embeds into $\langle I(\mathcal{B}) \rangle$ and $\langle F(\mathcal{A}) \rangle$ embeds into $\langle F(\mathcal{B}) \rangle$, iff $T(\mathcal{A}) \leq T(\mathcal{B})$. \square

Lemma 12. *Let \mathcal{A} and \mathcal{B} be two DNMG-chains. Then $\mathbf{V}(\mathcal{A}) \subseteq \mathbf{V}(\mathcal{B})$ iff $T(\mathcal{A}) \leq T(\mathcal{B})$. As a consequence, $\mathbf{V}(\mathcal{A}) = \mathbf{V}(\mathcal{B})$ iff $T(\mathcal{A}) = T(\mathcal{B})$.*

Proof. Assume first that $\mathbf{V}(\mathcal{A}) \subseteq \mathbf{V}(\mathcal{B})$. Then each chain in $\mathbf{V}(\mathcal{A})$ belongs to $\mathbf{V}(\mathcal{B})$, too. By Lemma 6, each finite chain $\mathcal{C} \in \mathbf{V}(\mathcal{A})$ embeds into both \mathcal{A} and \mathcal{B} . By Lemma 11 this occurs only if $T(\mathcal{C}) \leq T(\mathcal{A})$ and $T(\mathcal{C}) \leq T(\mathcal{B})$. But then $\langle X(\mathcal{A}) \rangle$ embeds into $\langle X(\mathcal{B}) \rangle$ for each $X \in \{F, I, S\}$, or they are both of infinite cardinality. In both cases $T(\mathcal{A}) \leq T(\mathcal{B})$.

For the other way round, assume $T(\mathcal{A}) \leq T(\mathcal{B})$. Take now any finite chain $\mathcal{C} \in \mathbf{V}(\mathcal{A})$. Then by Lemma 6, \mathcal{C} embeds into \mathcal{A} , and by Lemma 11, $T(\mathcal{C}) \leq T(\mathcal{A})$. By our standing assumption, $T(\mathcal{C}) \leq T(\mathcal{B})$, too. Whence, again by Lemma 11, \mathcal{C} embeds into \mathcal{B} , which implies that each finite chain in $\mathbf{V}(\mathcal{A})$ belongs to $\mathbf{V}(\mathcal{B})$, too. We conclude $\mathbf{V}(\mathcal{A}) \subseteq \mathbf{V}(\mathcal{B})$. \square

Notice that by Lemmas 2 and 8 and Theorem 2, if C is a set of DNMG-chains containing an infinite chain \mathcal{A} of any cardinality, then $\mathbf{V}(C) = \mathbf{V}(C')$ for $C' = (C \cup \{\mathcal{A}'\}) \setminus \{\mathcal{A}\}$, for a suitable \mathcal{A}' . More precisely, \mathcal{A}' is obtained from \mathcal{A} by replacing, for every $X \in \{F, I, S\}$ such that $X(\mathcal{A})$ is infinite, the subalgebra

$\langle X(A) \rangle$ with a fresh copy of a fixed denumerable $L(X)$ -chain \mathcal{B} , where we safely identify the extremes of \mathcal{B} with those of \mathcal{A} . We call \mathcal{A}' a *regular* chain.

Then we may assume that each subvariety of \mathbf{DNMG} is generated by a set of regular chains. Two regular \mathbf{DNMG} -chains \mathcal{A} and \mathcal{B} are isomorphic iff $T(\mathcal{A}) = T(\mathcal{B})$. We now fix, once and for all, one *representative* regular chain \mathcal{A}_T , for each triple $T \in (\omega + 1)^{(3)}$. Notice that given two regular representative \mathbf{DNMG} -chains \mathcal{A} and \mathcal{B} , we have that \mathcal{A} embeds into \mathcal{B} iff $T(\mathcal{A}) \leq T(\mathcal{B})$. A set C of \mathbf{DNMG} -chains is *irredundant* if each $\mathcal{A} \in C$ is representative regular and $\mathbf{V}(C \setminus \{\mathcal{A}\}) \subsetneq \mathbf{V}(C)$ for all $\mathcal{A} \in C$. Otherwise C is *redundant*.

Lemma 13. *Every irredundant set C of \mathbf{DNMG} -chains is finite. Moreover, the map $\mathcal{A} \mapsto T(\mathcal{A})$ is a bijection between C and $C_T = \{T(\mathcal{A}) : \mathcal{A} \in C\}$, which is the underlying set of a finite antichain of $(\omega + 1)^{(3)}$.*

Proof. Let C be an irredundant set of \mathbf{DNMG} -chains. Then, for every $\mathcal{A}, \mathcal{B} \in C$, with $\mathcal{A} \neq \mathcal{B}$ we have that $T(\mathcal{A})$ is incomparable with $T(\mathcal{B})$. Indeed, if not, by Lemma 12 either $\mathbf{V}(C \setminus \{\mathcal{A}\}) = \mathbf{V}(C)$ or $\mathbf{V}(C \setminus \{\mathcal{B}\}) = \mathbf{V}(C)$. In both cases we have a contradiction. By the previous observations we have that C_T must be the underlying set of an antichain of $(\omega + 1)^{(3)}$, and by Lemma 10 C_T is finite. The proof is settled by noticing that the map $\mathcal{A} \mapsto T(\mathcal{A})$ is a bijection between C and C_T . \square

Lemma 14. *The lattice $\Lambda(\mathbf{DNMG})$ of all subvarieties of \mathbf{DNMG} , ordered by inclusion, is isomorphic with the lattice $\Gamma(\mathbf{DNMG})$ of all irredundant sets of \mathbf{DNMG} -chains, ordered by inclusion.*

Proof. Clearly, each irredundant set of \mathbf{DNMG} -chains generates a subvariety of \mathbf{DNMG} , and each subvariety is generated by some irredundant set of \mathbf{DNMG} -chains. We prove that no subvariety can be generated by two distinct irredundant sets C, D of \mathbf{DNMG} -chains, with $C \neq D$: by Lemma 13 we can assume $C = \{\mathcal{A}_1, \dots, \mathcal{A}_h\}$ and $D = \{\mathcal{B}_1, \dots, \mathcal{B}_k\}$. By contradiction, suppose $\mathbf{V}(C) = \mathbf{V}(D)$.

Then, without loss of generality, there is a chain $\mathcal{A}_r \in C \setminus D$. By [6, Theorem 7], we have that the class of chains in $\mathbf{V}(C) = \mathbf{V}(D)$ coincides with the one in $\bigcup_{i=1}^h \mathbf{V}(\mathcal{A}_i)$ and with the one in $\bigcup_{i=1}^k \mathbf{V}(\mathcal{B}_i)$. This means that there is $\mathcal{B}_s \in D$ such that $\mathbf{V}(\mathcal{A}_r) \subseteq \mathbf{V}(\mathcal{B}_s)$, and clearly $\mathcal{B}_s \notin C$ (otherwise C would be redundant). This implies $T(\mathcal{A}_r) \leq T(\mathcal{B}_s)$ and hence \mathcal{A}_r embeds into \mathcal{B}_s , being both chains regular representative. On the other hand, for the same reasons, this in turn implies that \mathcal{B}_s embeds into some chain $\mathcal{A}_t \in C$. Note that $t \neq r$, as otherwise $\mathcal{A}_r \simeq \mathcal{B}_s$, and since both chains are regular representative we would conclude $\mathcal{A}_r = \mathcal{B}_s$, in contrast with the fact that $\mathcal{A}_r \notin D$. But then also \mathcal{A}_r embeds into \mathcal{A}_t , contradicting the fact that C is irredundant. It is obvious that both the bijective correspondence $C \mapsto \mathbf{V}(C)$ and its inverse are order-preserving. \square

Theorem 4. *The lattice $\Lambda(\mathbf{DNMG})$ is isomorphic with $\mathcal{AC}((\omega + 1)^{(3)})$.*

Proof. By Lemma 14, we prove that $\Gamma(\mathbf{DNMG})$ is isomorphic with $\mathcal{AC}((\omega + 1)^{(3)})$. Consider a set C of pairwise non-isomorphic representative chains in \mathbf{DNMG}

such that the set $C_T = \{T(\mathcal{A}) \mid \mathcal{A} \in C\}$ is not an antichain. Then there are chains \mathcal{A} and $\mathcal{B} \in C$ such that $T(\mathcal{A}) \leq T(\mathcal{B})$. By Lemma 12, $\mathbf{V}(\mathcal{A}) \subseteq \mathbf{V}(\mathcal{B})$. Whence, $\mathbf{V}(C) = \mathbf{V}(C \setminus \{\mathcal{A}\})$. That is C is redundant. Whence, each element of $\Gamma(\mathbb{DNMG})$ is a set C of chains such that C_T is an antichain, that is T maps $\Gamma(\mathbb{DNMG})$ to $\mathcal{AC}((\omega+1)^{(3)})$. T is injective by construction. For surjectivity, let $A \in \mathcal{AC}((\omega+1)^{(3)})$. With each triple $(a, b, c) \in A$ we associate the representative regular DNMG-chain $\mathcal{A}_{(a,b,c)}$ such that $T(\mathcal{A}_{(a,b,c)}) = (a, b, c)$. Call $C(A)$ the set of chains so obtained. Clearly, $T(C(A)) = A$ and all chains in $C(A)$ are representative and regular. Finally, for any triple $(a, b, c) \in A$ we have that $\mathbf{V}(C(A) \setminus \{\mathcal{A}_{(a,b,c)}\}) \subsetneq \mathbf{V}(C(A))$, for otherwise there is a triple $(d, e, f) \in A$ such that $(a, b, c) \leq (d, e, f)$, contradicting the fact that A is an antichain. Whence $C(A)$ is irredundant, that is, it belongs to $\Gamma(\mathbb{DNMG})$. Surjectivity is proved. \square

Corollary 2. *There are countably many subvarieties of \mathbb{DNMG} . Every subvariety of \mathbb{DNMG} is generated by a finite number of DNMG-chains.*

Proof. By Theorem 4, Lemma 10 and Corollary 1. \square

5 Uniform Axiomatisations

In this section we axiomatise all the subvarieties of \mathbb{DNMG} in a uniform way.

Definition 6. *For each integer $n > 0$ let Q_n be the term:*

$$Q_n = \bigsqcup_{1 \leq i \neq j \leq n+1} (x_i \Leftrightarrow x_j).$$

Furthermore let $Q_0 = 0$ and $Q_\omega = 1$, and $e_{\omega+1}(X) = 1$ for each $X \in \{F, I, S\}$.

Given $(a, b, c) \in (\omega+1)^{(3)}$, we write $\mathbf{V}(a, b, c)$ to mean the variety generated by a DNMG-chain \mathcal{A} such that $T(\mathcal{A}) = (a, b, c)$.

Theorem 5. *For each $(a, b, c) \in (\omega+1)^{(3)}$, the variety $\mathbf{V}(a, b, c)$ is the subvariety of \mathbb{DNMG} satisfying the following equation.*

$$(Q_a \sqcup e_{a+1}(S)) \sqcap (Q_b \sqcup e_{b+1}(I)) \sqcap (Q_c \sqcup e_{c+1}(F)) = 1.$$

Proof. Let \mathcal{D} be a DNMG-chain such that $T(\mathcal{D}) = (a, b, c)$. We show that $\mathcal{D} \models Q_a \sqcup e_{a+1}(S) = 1$. As a matter of fact, if $a = 0$ then by Lemma 3 $1 = e_{a+1}(d_1, \dots, d_{a+1}) = (Q_a \sqcup e_{a+1}(S))(d_1, \dots, d_{a+1})$, since $d_i \in D \setminus S(D)$ for all $i \in \{1, 2, \dots, a+1\}$. If $a = \omega$ then $(Q_a \sqcup e_{a+1}(S))(d_1, \dots, d_{m+1}) = 1$ for $Q_a = 1$. If a is a positive integer, then $(Q_a \sqcup e_{a+1}(S))(d_1, \dots, d_{a+1}) = 1$. Indeed, if $d_i \in S(D)$ for all $i \in \{1, 2, \dots, a+1\}$ then there are distinct indices i, j such that $d_i \Leftrightarrow d_j = 1$, for $|S(\mathcal{A})| = a$, whence $(Q_a \sqcup e_{a+1}(S))(d_1, \dots, d_{a+1}) = Q_a(d_1, \dots, d_{a+1}) = 1$. If on the other hand there is some $d_i \in D \setminus S(D)$, then, by Lemma 3, $e_{a+1}(S)(d_1, \dots, d_{a+1}) = 1 = (Q_a \sqcup e_{a+1}(S))(d_1, \dots, d_{a+1})$. We proceed similarly, to show that $\mathcal{D} \models Q_b \sqcup e_{b+1}(I) = 1$ and $\mathcal{D} \models Q_c \sqcup e_{c+1}(F) = 1$, *mutatis mutandis*. Let now \mathcal{D} be a DNMG-chain not in $\mathbf{V}(a, b, c)$, whence $T(\mathcal{D}) = (a', b', c')$

with $(a', b', c') \not\leq (a, b, c)$. If $a' > a$ then picking pairwise distinct elements $a_1, \dots, a_{a+1} \in S(D)$ we have that $Q_a(d_1, \dots, d_{a+1}) < 1$. Further, by Lemma 3, $e_{a+1}(S)(d_1, \dots, d_{a+1}) < 1$, too. We conclude that $\mathcal{D} \not\models Q_a \sqcup e_{a+1}(S) = 1$. The cases $b' > b$ and $c' > c$ are dealt with in the same manner, *mutatis mutandis*. \square

Clearly, if some index a, b, c is zero, the associated conjunct in the axiomatising equation can be simplified: $Q_a \sqcup e_{a+1}(S)$ to $e_{a+1}(S)$ and so forth. Analogously, if some index a, b, c is ω , the associated conjunct can be totally disregarded. For instance $\mathbf{V}(\omega, \omega, \omega) = \mathbb{DNMG}$, and as a matter of fact the axiomatising equation in this case is identically $1 = 1$. As \mathbb{DNMG} contains all major subvarieties of \mathbb{WNM} , several varieties of the form $\mathbf{V}(a, b, c)$ have already been studied and axiomatised in the literature. The following theorems report on this aspect.

Theorem 6. $\mathbf{V}(0, 0, 0) = \mathbb{B}$. For each integer $n > 0$, $\mathbf{V}(n, 0, 0) = \mathbb{G}_{n+2}$, $\mathbf{V}(0, n, 0) = \mathbb{NM}_{2n+2}$, $\mathbf{V}(0, 0, n) = \mathbb{DP}_{n+2}$. Also, $\mathbf{V}(\omega, 0, 0) = \mathbb{G}$, $\mathbf{V}(0, \omega, 0) = \mathbb{NM}^-$ and $\mathbf{V}(0, 0, \omega) = \mathbb{DP}$.

Proof. Immediate by [6, 17], and Definition 5. \square

We now show how the axiomatisation provided in Theorem 5 can be simplified, when exactly one element in the triplet (a, b, c) is zero. In this case $\mathbf{V}(a, b, c)$ is either a subvariety of $\mathbb{RDP} = \mathbf{V}(\omega, 0, \omega)$, or of one of the following two subvarieties. The variety $\mathbb{DNM} = \mathbf{V}(0, \omega, \omega)$, axiomatised as \mathbb{WNM} plus:

$$(\sim \sim x \Rightarrow x) \sqcup (\sim \sim x \Leftrightarrow \sim x) = 1. \quad (\text{IF})$$

The variety $\mathbb{NMG}^- = \mathbf{V}(\omega, \omega, 0)$, axiomatised as \mathbb{NMG} plus:

$$\sim((\sim(x^2))^2) \Leftrightarrow (\sim((\sim x)^2))^2 = 1. \quad (\text{NF})$$

Theorem 7. 1. For all $b, c \in (\omega + 1)^{(2)}$ with $b \neq 0 \neq c$, the variety $\mathbf{V}(0, b, c)$ is axiomatised as \mathbb{DNM} plus

$$(Q_b \sqcup e_{b+1}(I)) \sqcap (Q_c \sqcup e_{c+1}(F)) = 1.$$

2. For all $a, c \in (\omega + 1)^{(2)}$ with $a \neq 0 \neq c$, the variety $\mathbf{V}(a, 0, c)$ is axiomatised as \mathbb{RDP} plus

$$(Q_a \sqcup e_{a+1}(S)) \sqcap (Q_c \sqcup e_{c+1}(F)) = 1.$$

3. For all $a, b \in (\omega + 1)^{(2)}$ with $b \neq 0 \neq a$, the variety $\mathbf{V}(a, b, 0)$ is axiomatised as \mathbb{NMG}^- plus

$$(Q_a \sqcup e_{a+1}(S)) \sqcap (Q_b \sqcup e_{b+1}(I)) = 1.$$

Proof. Immediate by Theorem 5.

We now provide the general criterion to axiomatise the subvarieties of \mathbb{DNMG} .

Theorem 8. *Let $C = \{\mathcal{A}_i\}_{i \in I}$ be an irredundant set of DNMG-chains. Let further $t_i(x_1, \dots, x_{n_i}) = 1$ be the equation axiomatising $\mathbf{V}(\mathcal{A}_i)$ for each $i \in I$, as given by Theorem 5. Then $\mathbf{V}(C)$ contains exactly the DNMG-algebras satisfying the equation*

$$\bigsqcup_{i \in I} t_i(y_{i,1}, \dots, y_{i,n_i}) = 1,$$

where all the variables $y_{i,j}$, for $i \in I$, and $j \in \{1, \dots, n_i\}$, are pairwise distinct.

Proof. First notice that by Corollary 2, $\bigsqcup_{i \in I} t_i = 1$ is indeed an equation, as I is a finite index set. The proof is settled by noting that $\mathbf{V}(C) = \bigsqcup_{i \in I} \mathbf{V}(\mathcal{A}_i)$, and by using [15, Lemma 5.25]. \square

Corollary 3. *Every element of $\Lambda(\text{DNMG})$ is the join of a finite set of join irreducible elements.*

Proof. By [2, Theorem 5.1] a variety of MTL-algebras is join irreducible, in the lattice of the subvarieties of MTL, if and only if it is generated by a single chain. The claim follows by Theorem 8. \square

Theorem 9. *DNMG is the smallest subvariety in $\Lambda(\text{DNMG})$ which contains DP , NM^- , \mathbb{G} and it is generated by a single chain.*

Proof. Immediate by Theorem 6 and Lemma 12, since $\text{DNMG} = \mathbf{V}(\omega, \omega, \omega)$. \square

Remark 1. Notice that $\text{NM} = \mathbf{V}(0, \omega, 1)$, and its lattice of subvarieties is given by all antichains $C \in \mathcal{AC}((\omega + 1)^{(3)})$ such that all $T \in C$ have either the form $T = (0, b, 1)$ or $T = (0, b, 0)$ for some integer $b \geq 0$, or $T = (0, \omega, 0)$, whose corresponding variety is NM^- .

The almost minimal subvarieties of DNMG are exactly $\mathbb{G}_3 = \mathbf{V}(1, 0, 0)$, $\text{NM}_4 = \mathbf{V}(0, 1, 0)$, and $\text{NM}_3 = \text{DP}_3 = \mathbf{V}(0, 0, 1)$. Whence they coincide with the almost minimal subvarieties of WNM (see [3]).

By Lemma 7 and [10, Theorem 3.5], every variety of DNMG-algebras of the form $\mathbf{V}(a, b, c)$ is such that the corresponding logic has the strong single chain completeness (see [2, 22]).

The subvarieties of DNMG generated by a standard algebra are exactly $\mathbb{G} = \mathbf{V}(\omega, 0, 0)$, $\text{NM} = \mathbf{V}(0, \omega, 1)$, $\text{NMG} = \mathbf{V}(\omega, \omega, 1)$, $\text{RDP} = \mathbf{V}(\omega, 0, \omega)$, $\text{DNM} = \mathbf{V}(0, \omega, \omega)$, and, clearly, $\text{DNMG} = \mathbf{V}(\omega, \omega, \omega)$.

Finally, $\text{EMTL} = \mathbf{V}(\{(\omega, 0, 0), (0, 0, \omega)\})$ is an example of a subvariety of DNMG which cannot be generated by a single chain [6].

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