

Asynchronous ASMs

Phase Synchronization on Undirected Trees

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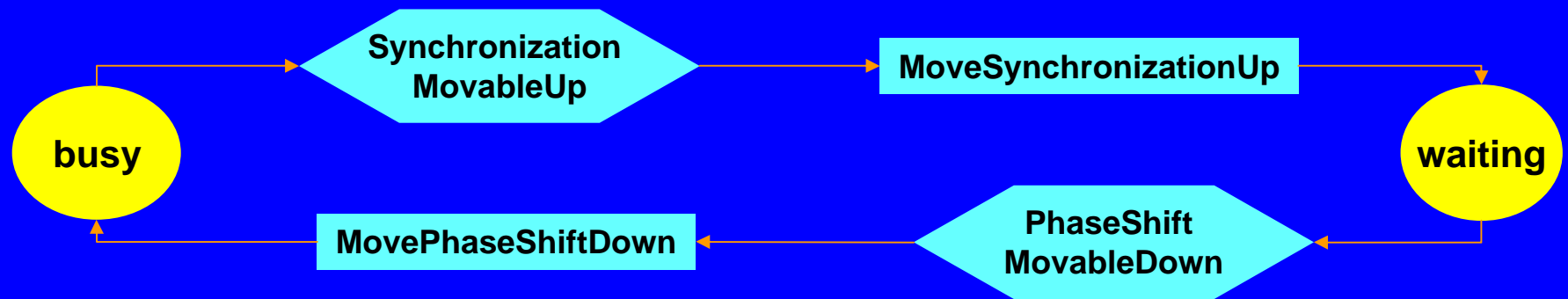
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Phase Synchronization : problem statement

- Goal: Design a distributed algorithm that guarantees synchronized execution of phases of computations, occurring at nodes of an undirected tree (connected acyclic network), using only communication between tree neighbors
- Algorithmic Idea: every agent (tree node) before becoming busy in a new phase has to
 - synchronize with its neighbor nodes, moving the synchronization up along the tree structure
 - wait until all agents are waiting for the phase increase to start
 - increase its phase - when the phase shift becomes movable down - and move the phase shift further down along the tree structure
 - until all nodes have become busy in the new phase

Phase Synchronization ASM: Agent Signature

- **Agent** : an undirected tree of agents equipped with:
 - **neighb** \subseteq Agent (external function)
 - **phase**: the current phase
 - **synchPartner**, **waitPartner**: Agent yields the neighbor node an agent is synchronized with (in moving the synchronization up or down the tree) resp. waiting for (to reverse the previous upward synchronization downwards)
 - **ctl_state** : {**busy**, **waiting** }
- **Initially**
 - **ctl_state** = **busy**, **phase** = 0, **synchPartner** = **waitPartner** = **undef**



SynchronizationMovableUp $\circ \exists y \in \text{neighb} \quad \text{SynchronizationMovableUpTo}(y)$
 for some neighb y all other neighbors synchronize with self in phase(self)

MoveSynchronizationUp $\circ \text{choose } y \in \text{neighb} \text{ with } \text{SynchronizationMovableUpTo}(y)$
 synchronize self with y in phase(self) $\text{MoveSynchronizationUpTo}(y)$

SynchronizationMovableUpTo(y) $\circ \forall z \in \text{neighb} - \{y\}$
 $z.\text{synchPartner}(\text{phase}(\text{self})) = \text{self}$

PhaseShiftMovableDown
 waitPartner synchronizes
 with self in phase(self)

MoveSynchronizationUpTo(y) $\circ \forall z \in \text{neighb} - \{y\}$
 $z.\text{synchPartner}(\text{phase}(\text{self})) := \text{undef}$
 $\text{self}.\text{synchPartner}(\text{phase}(\text{self})) := y, \quad \text{self}.\text{waitPartner}(\text{phase}(\text{self})) := y$

MovePhaseShiftDown $\circ \quad \text{phase} := \text{phase} + 1$
 $\forall z \in \text{neighb} - \{\text{waitPartner}\} \quad \text{self}.\text{synchPartner}(\text{phase}(\text{self})) := z$
 $\text{waitPartner}.\text{synchPartner}(\text{phase}(\text{self})) := \text{undef}$

Phase Synchronization ASM : Correctness & Liveness

- Proposition: In every infinite distributed run of agents equipped with the ASM for phase synchronization:
 - in any state any two busy agents are in the same phase (correctness property)
 - if every enabled agent will eventually make a move, then each agent will eventually reach each phase (liveness property)

Phase Synch: **PhaseShiftMovableDown Lemma**

- Lemma: For every phase p , whenever in a run a state is reached where for the first time
 $u.\text{waitPartner}(p)=v$ synchronizes with u in phase p ,
every element of
 $\text{subtree}(u,v) \cup \text{subtree}(v,u) \cup \{u,v\}$
is waiting in phase p
 where $\text{subtree}(x,y) = \{n \mid n \text{ reachable by a path from } x \text{ without touching } y\}$

Proof of PhaseShiftMovableDown Lemma

- Proof of the lemma follows from the following two claims.
- Claim 1. When Synchronization is moved up in $\text{phase}(u)=p$ from busy u to $v=u.\text{synchPartner}(p)$, the elements of $\text{subtree}(u,v)$ are waiting, u becomes waiting, and they all remain so until the next application of a Shift rule “MovePhaseShiftDown”.
 - Claim 1 follows by induction on the applications of the Synchronization rule.
 - $n=1$: true at the leaves
 - $n+1$: follows by induction hypothesis and Synchronization rule from
$$\text{subtree}(u,v) = \hat{E}_{i < n} \text{subtree}(u_i, u) \cup \{u_0, \dots, u_{n-1}\} \text{ for } \text{neighb}(u) = \{u_0, \dots, u_n\}$$
with $v = u_n$ (by connectedness)

Proof of PhaseShiftMovableDown Lemma

- Claim 2. For every infinite run and every phase p , a state is reached in which all agents are waiting in phase p and for some agent u its waitPartner in phase p synchronizes with u in phase p .
 - Follows by induction on p .

Reference

- W.Reisig: Elements of Distributed Algorithms
Springer-Verlag 1998
 - See Section 36 (in particular Fig. 36.1) and a correctness and liveness proof in Section 81
- E. Börger, R. Stärk: Abstract State Machines. A Method for High-Level System Design and Analysis
Springer-Verlag 2003, see <http://www.di.unipi.it/AsmBook>
 - See Chapter 6.1